



# Advanced Functional Programming

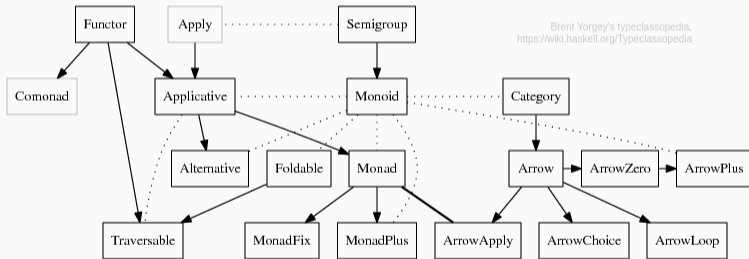
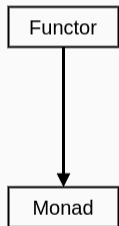
Applicative, Traversable, and Foldable

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## How to learn Haskell



Brent Yorgey's typeclassopedia,  
<https://wiki.haskell.org/Typeclassopedia>

1. learn monads

2. learn the rest of the typeclassopedia

## Sequencing IO operations

Some realistic Haskell code:

```
main = do
```

```
  fs <- getArgs
```

```
  cs <- mapIO readFile fs
```

```
  putStrLn (concat cs)
```

```
mapIO :: (a -> IO b) -> [a] -> IO [b]
```

```
mapIO f []      = return []
```

```
mapIO f (c : cs) = do x <- f c
                      xs <- mapIO f cs
                      return (x : xs)
```

- What does `main` do?
- What's wrong with this code?

## Code style

Which function would you rather review?

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (c : cs) = let x = f c
                  xs = map f cs
                  in (x : xs)
```

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (c : cs) = f c : map f cs
```

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map f (c : cs) = let x = f c
                  xs = map f cs
                  in (x : xs)
```

- Verbosity!
- Local variables!

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (c : cs) = f c : map f cs
```

## Code style

Which function would you rather review?

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (c : cs) = let x = f c
                  xs = map f cs
                  in (x : xs)
```

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (c : cs) = f x : map f xs
```

- Verbosity!
- Local variables!  
*There are only two hard things in Computer Science: cache invalidation and naming things.*  
– Phil Karlton

## Code style

Which function would you rather review?

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (c : cs) = let x = f c
                  xs = map f cs
                  in (x : xs)
```

```
map :: (a -> b) -> [a] -> [b]
map f []          = []
map f (c : cs) = f x : map f xs
```

- Verbosity!
- Local variables!  
*There are only two hard things in Computer Science: cache invalidation and naming things.*  
– Phil Karlton
- do notation forces un-idiomatic style?

## Using liftM2

The liftM2 function is defined as follows:

```
liftM2 :: (a -> b -> c) -> m a -> m b -> m c
liftM2 f ma mb = do a <- ma
                   b <- mb
                   return (f a b)
```

Using liftM2 we can write:

```
mapIO :: (a -> IO b) -> [a] -> IO [b]
mapIO f []      = return []
mapIO f (c:cs) = liftM2 (:) (f c) (mapIO f cs)
```

This even works for *any monad*, not just the IO monad.

## More lifting functions

- `liftM` (or `fmap`) lifts functions `a -> b`
- `liftM2` lifts functions `a -> b -> c`
- ...
- `liftM5` lifts functions `a -> b -> c -> d -> e -> f`

Do we need a `liftMn` for every `n`?

**Time to derive liftMn!**

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## Using ap

The ap function is defined as follows:

```
ap :: Monad m => m (a -> b) -> m a -> m b
ap mf mx = do f <- mf
              x <- mx
              return (f x)
```

Using ap we can write:

```
mapIO :: (a -> IO b) -> [a] -> IO [a]
mapIO f []      = return []
mapIO f (c:cs) =
  return (:) `ap` f c `ap` mapIO f cs
```

## Evaluating expressions

Another example:

```
data Expr v = Var v | Val Int | Add (Expr v) (Expr v)
```

```
type Env v = Map v Int
```

```
eval :: Expr v -> Env v -> Int
```

```
eval (Var v) env = lookup v env
```

```
eval (Val i) env = i
```

```
eval (Add l r) env = (eval l env) + (eval r env)
```

Once again, we are passing around an environment that is only really used in the Var branch.

## An applicative alternative

```
const :: a -> (env -> a)
```

```
const x = \env -> x
```

```
s :: (env -> a -> b) -> (env -> a) -> (env -> b)
```

```
s ef es env = (ef env) (es env)
```

```
eval :: Expr v -> Env v -> Int
```

```
eval (Var v) = lookup v
```

```
eval (Val i) = const i
```

```
eval (Add l r) =
```

```
  const (+) `s` (eval l) `s` (eval r)
```

The `s` combinator lets us 'apply' one computation expecting an environment to another.

## Transposing matrices

```
transpose :: [[a]] -> [[a]]
transpose []          = repeat []
transpose (xs : xss) = zipWith (:) xs (transpose xss)
```

Can we play the same trick and find a combinator that will 'apply' a list of functions to a list of arguments?

```
zapp :: [a -> b] -> [a] -> [b]
zapp (f : fs) (x : xs) = (f x) : (zapp fs xs)
```

```
transpose (xs : xss) =
  repeat (:) `zapp` xs `zapp` transpose xss
```

## What is the pattern?

What do these functions have in common?

```
ap    :: IO (a -> b)    -> IO a        -> IO b
```

```
s     :: (env -> a -> b) -> (env -> a) -> (env -> b)
```

```
zapp :: [a -> b]        -> [a]         -> [b]
```

## Applicative (**applicative functors**)

```
class Functor f => Applicative f where
  pure  :: a -> f a
  (<*>) :: f (a -> b) -> f a -> f b
```

Note that Functor is a superclass of Applicative.

We can also define fmap in terms of the applicative operations:

```
(<$>) :: Functor f => (a -> b) -> f a -> f b
```

## Applicative (**applicative functors**)

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class Functor f => Applicative f where
  pure  :: a -> f a
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```

Note that Functor is a superclass of Applicative.

We can also define fmap in terms of the applicative operations:

```
(<$>) :: Functor f => (a -> b) -> f a -> f b
```

```
(<$>) f fx = pure f <*> fx
```

## Using Applicative operators

This type class leads to a certain code style:

```
mapIO :: (a -> IO b) -> [a] -> IO [b]
```

```
mapIO f [] = return []
```

```
mapIO f (c:cs) = (:) <$> f c <*> mapIO f cs
```

## Relating Applicative **functors** and Monads

- Every monad can be given an applicative functor interface.

```
instance Monad m => Applicative m where
  pure :: a -> m a
  pure = return
  mf <*> mx = do f <- mf; x <- mx; return (f x)
```

- But this may not always be the 'right' choice. For example, we have seen the 'zapp' applicative instance for lists; using the instance arising from the list monad gives very different behaviour!
- But not every applicative functor is a monad...

## Monads vs. applicative functors (1)

```
( $\$$ )    ::          (a -> b) -> a -> b
(< $\$$ >)  :: Functor f  => (a -> b) -> f a -> f b
(<*>)  :: Applicative f => f (a -> b) -> f a -> f b
(=<<)  :: Monad f     => (a -> f b) -> f a -> f b
```

- All generalize function application
- With (< $\$$ >), only one arg has effects
- With (<\*>), both args have effects, independently of each other
- With (=<<), subsequent actions can depend on the results of effects: depending on the character the user enters, respond differently.

## Monads vs applicative functors (2)

- There are more Applicative functors than there are monads; but monads are more powerful!
- If you *have* an Applicative functor, that's good; having a monad is better.
- If you *need* a monad, that's good; only needing an Applicative functor is better.
- With applicative functors, the structure is statically determined (and can be analyzed or optimized from inside Haskell).

```
miffy :: Monad m => m Bool -> m a -> m a -> m a
```

## Imprecise but catchy slogans

*Monads are programmable semi-colons!*

*Applicatives are programmable function application!*

## Applicative functor laws

- identity

$$\text{pure id } \langle * \rangle u = u$$

- composition

$$\text{pure } (.) \langle * \rangle u \langle * \rangle v \langle * \rangle w = u \langle * \rangle (v \langle * \rangle w)$$

- homomorphism

$$\text{pure } f \langle * \rangle \text{pure } x = \text{pure } (f x)$$

- interchange

$$u \langle * \rangle \text{pure } x = \text{pure } (\backslash f \rightarrow f x) \langle * \rangle u$$

## Spot the pattern

```
mapIO :: (a -> IO b) -> [a] -> IO [b]
mapIO f []      = return []
mapIO f (c:cs) = (:) <$> f c <*> mapIO f cs
```

```
mapEither :: (a -> Either err b) -> [a] -> Either err [b]
mapEither f [] = return []
mapEither f (c:cs) = (:) <$> f c <*> mapEither f cs
```

Both these functions map an effectful function over a list, and collect the results.

## Traversing lists

We can define a new function to capture this pattern:

```
mapM :: Monad m => (a -> m b) -> [a] -> m [b]
```

```
mapM _ [] = return []
```

```
mapM f (x:xs) = (:) <$> f x <*> mapM f xs
```

But the Monad constraint is quite restrictive: we don't use ( $\gg=$ ), and some types don't have it...

## Traversing lists without Monads

Since base 4.3.0.0, `mapM = traverse`.

```
traverse :: Applicative f => (a -> f b) -> [a] -> f [b]
```

```
traverse _ [] = pure []
```

```
traverse f (x:xs) = (:) <$> f x <*> traverse f xs
```

Clearly we can traverse lists in this fashion – but what other data types support such an operation?

The Traversable class captures those types that can be traversed in this fashion:

```
class Traversable t where
  traverse  :: Applicative f
            => (a -> f b) -> t a -> f (t b)
  sequenceA :: Applicative f
            => t (f a) -> f (t a)
```

The Traversable class captures those types that can be traversed in this fashion:

```
class Traversable t where
  traverse  :: Applicative f
            => (a -> f b) -> t a -> f (t b)
  sequenceA :: Applicative f
            => t (f a) -> f (t a)
```

## Question:

Define traverse and sequenceA in terms of each other

## Traversable with defaults

```
class Traversable t where

  traverse  :: Applicative f
            => (a -> f b) -> t a -> f (t b)
  traverse f = sequenceA . fmap f

  sequenceA :: Applicative f
            => t (f a) -> f (t a)
  sequenceA = traverse id
```

## Traversable: example

```
data Expr v = Var v | Val Int | Add (Expr v) (Expr v)
```

```
instance Traversable Expr where
```

```
  traverse :: Applicative f => (a -> f b) => Expr a -> f (Expr b)
```

```
  traverse f (Var v)   = Var <$> f v
```

```
  traverse f (Val x)   = pure (Val x)
```

```
  traverse f (Add l r) =
```

```
    Add <$> traverse f l <*> traverse f r
```

In general, `traverse` is just like `fmap` – only in applicative style.

## Introducing Foldable

In the Haskell libraries, Traversable is defined slightly differently.

```
class (Functor t, Foldable t)
    => Traversable t where
```

What is the Foldable class?

## Folding a list

The `foldr` function on lists captures a common pattern – think of it as a functional for-loop.

```
foldr :: (a -> b -> b) -> b -> [a] -> b
```

```
foldr f y [] = y
```

```
foldr f y (x:xs) = f x (foldr f y xs)
```

We can use it to define all kinds of simple list traversals:

```
sum = foldr (+) 0
```

```
maximum = foldr max minBound xs
```

```
(++) = \xs ys -> foldr (:) ys xs
```

```
concat = foldr (++) []
```

```
map = \f -> foldr (\x xs -> f x : xs) []
```

## Folding: beyond lists

There are many other data types that support some form of fold operator.

```
data Tree a = Empty
            | Leaf a
            | Node (Tree a) (Tree a)

foldTree :: (a -> b -> b) -> b -> Tree a -> b
foldTree f y Empty      = y
foldTree f y (Leaf x)   = f x y
foldTree f y (Node l r) =
    foldTree f (foldTree f y r) l
```

Note that *generic programming* gives a slightly more precise account.

# Foldable

```
class Foldable f where
  foldr  :: (a -> b -> b) -> b -> f a -> b
  foldMap :: Monoid m => (a -> m) -> f a -> m
```

Sometimes it can be easier to define the foldMap function – but what is a Monoid?

## Intermezzo: monoids

```
class Monoid a where  
  mempty  :: a  
  mappend :: a -> a -> a
```

Here mempty should be the unit of the associative operator mappend.

## Monoids everywhere

- Bool using && and True
- Bool using or and False
- Int using + and 0
- Int using \* and 1
- Int using max and minBound
- List a using ++ and []
- Imperative programs using ; and skip
- a -> a using . and id
- ...

Monoids pop up everywhere!

## Defining foldMap

Instead of defining fold, sometimes it can be easier to define foldMap:

```
foldMap :: Monoid m => (a -> m) -> Tree a -> m
foldMap f Empty      = mempty
foldMap f (Leaf x)   = f x
foldMap f (Node l r) =
  foldMap f l `mappend` foldMap f r
```

You need to apply `f` to all the `a` values in the tree and combine subtrees using `mappend`.

# Why?

What is the point of all this abstraction?

We all agree (I hope!) that `foldr` is useful for lists.

```
sum      = foldr (+) 0
```

```
maximum = foldr max minBound xs
```

```
(++) ys = foldr (:) ys xs
```

```
concat  = foldr (++) []
```

```
map f   = foldr (\x xs -> f x : xs) []
```

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(++) ys = foldr (:) ys xs
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```
concat  = foldr (++) []
```

```
map f   = foldr (\x xs -> f x : xs) []
```

... but we can now give definitions that work for *any* foldable structure.

# Why?

What is the point of all this abstraction?

We all agree (I hope!) that `foldr` is useful for lists.

```
sum    :: Foldable f => f Int -> Int
```

```
sum    = foldr (+) 0
```

```
maximum :: Foldable f => f Int -> Int
```

```
maximum = foldr max minBound
```

```
flatten :: Foldable f => f a -> [a]
```

```
flatten = foldMap (\x -> [x]) (++)
```

As a slightly less trivial example, consider the any function:

```
any :: (a -> Bool) -> [a] -> Bool
```

```
any p []      = False
```

```
any p (x:xs) = p x || any p xs
```

How can we generalize this to work on *any* traversable type?

## Mighty Booleans

Let's start by finding the right monoidal structure.

Instead of defining an instance for `Bool`, introducing a new type can help clarify the monoidal structure we are using.

```
newtype Might = Might { might :: Bool }
```

```
instance Monoid Might where
```

```
  mempty = Might False
```

```
  (Might x) `mappend` (Might y) = Might (x || y)
```

## Generic any

```
any :: Foldable f => (a -> Bool) -> f a -> Bool  
any p = might . foldMap (Maybe . p)
```

Many other functions can be generalized similarly.

## The Foldable-Traversable Proposal

As of GHC 7.10, many Prelude functions have been generalized to work over *any* traversable structure – and not just lists.

Suppose we have a data type for binary trees, with the obvious traversable/foldable instances:

```
data Tree a where
  Leaf  :: a -> Tree a
  Node  :: Tree a -> Tree a -> Tree a
```

We can use the prelude functions we are used to over this data structure too.

## Example folds over trees

```
> let t = (Node (Leaf 1) (Node (Leaf 2) (Leaf 3)))
```

```
> any isEven t
```

```
True
```

```
> length t
```

```
3
```

```
> elem 3 t
```

```
True
```

We no longer need to define specialized functions for trees.

## Exercise

What is the foldable instance for `Maybe`?

What about the foldable instance for `pairs`?

But there are also quite surprising examples:

```
minimum (1,1000)
```

```
length (lookup 4 [(2,“Hello”), (4,“World”), (5,“!”)])
```

```
null (lookup 3 [])
```

## Drawbacks

But there are also quite surprising examples:

```
minimum (1,1000)
```

```
length (lookup 4 [(2,“Hello”), (4,“World”), (5,“!”)])
```

```
null (lookup 3 [])
```

Sometimes code may type check, where you would have liked to see a type error.

## Introducing arrows

If applicative functors generalize the notion of application, can we find a similar abstraction over functions and function composition?

Yes! There is more than a decade of work investigating functional programming using *Arrows*.

```
class Arrow a where
  arr    :: (b -> c) -> a b c
  (>>>) :: a b c -> a c d -> a b d
  first :: a b c -> a (b,d) (c,d)
```

- Just like applicative functors and monads, arrows have several associated laws.
- Many programs using arrows require additional operations – similar to classes such as `MonadPlus`.
- GHC supports special syntax for programming with Arrows, similar to the `do` notation for monads.

## Historical context

- Monads were originally studied in the context of program language semantics.
- Only later, was their importance for *structuring* programs discovered (and subsequently, modelling IO)
- Arrows (Hughes 2000) were proposed as an alternative to monads, but they have not been widely adopted.
- More recently, applicative functors have gained a lot of traction in the Haskell community (McBride and Paterson 2008), generalising the interface by Duponcheel and Swierstra (1996).
- Applicative functors, together with the associated Traversable and Foldable classes, are now part of the Haskell standard.

## Why care?

Functional programmers are addicted to abstraction: as soon as they spot a pattern, they typically want to abstract over it.

The type classes we have seen today, such as monads, applicative functors, foldable, and traversable, all capture some common pattern.

Using these patterns can save you some boilerplate code.

But *understanding* these patterns can help guide your design.

Is my type a monad? Or is it just applicative? Can I find a Traversable instance? Why not?

## Further reading

- *Applicative programming with effects*, McBride and Paterson
- *Monoids: Theme and Variations*, Brent Yorgey
- *Programming with arrows*, John Hughes
- *Idioms are oblivious, arrows are meticulous, monads are promiscuous*, Lindley, Wadler and Yallop
- and much much more!

## **Appendix: Time for type trickery!**

---

```
class (Functor t, Foldable t)
  => Traversable t where
  traverse  :: Applicative f
            => (a -> f b) -> t a -> f (t b)
  sequenceA :: Applicative f
            => t (f a) -> f (t a)
```

We can write `fmap` and `foldMap` using `traverse` or `sequenceA`

```
class (Functor t, Foldable t)
  => Traversable t where
  traverse  :: Applicative f
            => (a -> f b) -> t a -> f (t b)
  sequenceA :: Applicative f
            => t (f a) -> f (t a)
```

We can write `fmap` and `foldMap` using `traverse` or `sequenceA`

**Let's do it!**

```
newtype Id a = Id { getId :: a }
```

```
myFmap :: Traversable f => (a -> b) -> f a -> f b
```

```
myFmap f = getId . traverse (Id . f)
```

## foldMap **via** traverse

```
traverse :: Applicative f  
         => (a -> f b) -> t a -> f (t b)  
foldMap  :: Monoid m  
         => (a -> m)   -> t a -> m
```

## foldMap **via** traverse

```
traverse :: Applicative f
          => (a -> f b) -> t a -> f (t b)

foldMap  :: Monoid m
          => (a -> m)  -> t a -> m
```

You need to find an Applicative which behaves like a monoid

```
myFoldMap f = getMonoidApplicative
              . traverse (MonoidApplicative . f)
```

## foldMap **via** traverse **using** Const

The functor you need is called Const

```
data Const k a = Const { getConst :: k }
```

```
instance Functor (Const k) where  
  fmap _ (Const x) = Const x
```

What about the Applicative instance?

## foldMap **via** traverse **using** Const

```
instance Monoid k => Applicative (Const k) where
  pure = Const mempty
  Const x <*> Const y
    = Const (x `mappend` y)
```