

Parser Design



common pitfalls, and how to avoid them

Announcements

Lab 1: Group registration

- Ask Johan.

Lab 1: base = 4.19.2.0

- The autograder uses this.
- Bump at your own risk.

Werkcollege today: 5cp bonus

- Work on lab 1 or exercises.

Today: last parsing lecture before Lab 1



Previously...

String

e.g.

	<i>Groningen</i>	<i>8:37</i>
<i>10:15</i>	<i>Utrecht</i>	<i>10:21</i>
<i>11:05</i>	<i>Den Haag</i>	<i>11:23</i>
<i>11:34</i>	<i>Leiden</i>	

Previously...

Language

e.g.

| “the set of all well-formed travel
schedules”

Previously...

Grammar

e.g.

$S \rightarrow \textit{Start Wait}^* \textit{End}$

$\textit{Start} \rightarrow \textit{StationName Time}$

$\textit{Wait} \rightarrow \textit{Time StationName Time}$

$\textit{End} \rightarrow \textit{Time StationName}$

Previously...

Abstract Syntax

e.g.

```
data TSched = TSched
  { start ::      (String, Time)
  , waits :: [(Time, String, Time)]
  , end   ::      (Time, String)
  }
```

Previously...

Parsers

e.g.

```
parseTSched :: String  
            → [(TSched, String)]
```

Quiz: parser for travel schedules?

Groningen 8:37
10:15 Utrecht 10:21
11:05 Den Haag 11.23
11:34 Leiden

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```
{- 1 -}      many (pStart <|> pWait <|> pEnd)
{- 2 -} TSched <$> (many (pStart <|> pWait <|> pEnd))
{- 3 -} TSched <$> pStart <*>      pWaits <*> pEnd
{- 4 -} TSched <$> pStart <*> many pWait <*> pEnd
{- 5 -} TSched <$> pStart <$> many pWait <$> pEnd
{- 6 -} TSched <*> pStart <$> many pWait <*> pEnd
```

Not 1

`many :: ... → Parser c [a]`

```
--          many (pStart <|> pWait <|> pEnd)
{- 2 -} TSched <$> (many (pStart <|> pWait <|> pEnd))
{- 3 -} TSched <$> pStart <*> pWaits <*> pEnd
{- 4 -} TSched <$> pStart <*> many pWait <*> pEnd
{- 5 -} TSched <$> pStart <$> many pWait <$> pEnd
{- 6 -} TSched <*> pStart <$> many pWait <*> pEnd
```

Not 2

TSched :: ... → ... → ... → TSched

many :: ... → Parser c [a]

```
--          many (pStart <|> pWait <|> pEnd)
--      TSched <$> (many (pStart <|> pWait <|> pEnd))
{- 3 -} TSched <$> pStart <*> pWaits <*> pEnd
{- 4 -} TSched <$> pStart <*> many pWait <*> pEnd
{- 5 -} TSched <$> pStart <$> many pWait <$> pEnd
{- 6 -} TSched <*> pStart <$> many pWait <$> pEnd
```

Yes 3,4

TSched :: ... → ... → ... → TSched
many :: ... → Parser c [a]

```
--          many (pStart <|> pWait <|> pEnd)
--      TSched <$> (many (pStart <|> pWait <|> pEnd))
{- C -} TSched <$> pStart <*> pWaits <*> pEnd
{- D -} TSched <$> pStart <*> many pWait <*> pEnd
{- E -} TSched <$> pStart <$> many pWait <$> pEnd
{- F -} TSched <*> pStart <$> many pWait <$> pEnd
```

Not 5,6

```
TSched :: ... → ... → ... → TSched
infix1 (<$>) :: (a → b) → Parser c a → Parser c b
infix1 (<*>) :: Parser c (a → b) → Parser c a → Parser c b
infix1 (<$)  :: a → Parser c b → Parser c a
infix1 (<*)  :: Parser c a → Parser c b → Parser c a
```

```
--          many (pTStart <|> pTEnd <|> pTLine)
--      TSched <$> (many (pTStart <|> pTEnd <|> pTLine))
{- 3 -} TSched <$> pStart <*> pTLine <*> pWaits <*> pEnd
{- 4 -} TSched <$> pStart <*> pTLine <*> many pWait <*> pEnd
--      TSched <$> pStart <$> pTLine <$> many pWait <$> pEnd
--      TSched <*> pStart <$> pTLine <*> many pWait <$> pEnd
```

Combinator zoo

$(\langle \$ \rangle) :: (a \rightarrow b) \rightarrow \text{Parser } c \ a \rightarrow \text{Parser } c \ b$
 $(\langle * \rangle) :: \text{Parser } c \ (a \rightarrow b) \rightarrow \text{Parser } c \ a \rightarrow \text{Parser } c \ b$
 $(\langle \$ \rangle) :: a \rightarrow \text{Parser } c \ b \rightarrow \text{Parser } c \ a$
 $(\langle * \rangle) :: \text{Parser } c \ a \rightarrow \text{Parser } c \ b \rightarrow \text{Parser } c \ a$

Combinator zoo, abbreviated

$(\langle \$ \rangle) ::$	$(a \rightarrow b) \rightarrow$	$f \ a \rightarrow$	$f \ b$
$(\langle * \rangle) ::$	$f \ (a \rightarrow b) \rightarrow$	$f \ a \rightarrow$	$f \ b$
$(\langle \$ \rangle) ::$	$a \rightarrow$	$f \ b \rightarrow f \ a$	
$(\langle * \rangle) ::$	$f \ a \rightarrow$	$f \ b \rightarrow f \ a$	

Combinator zoo, grouped

$(\langle \$ \rangle) ::$	$(a \rightarrow b) \rightarrow$	$f \ a \rightarrow$	$f \ b$
$(\langle * \rangle) ::$	$f \ (a \rightarrow b) \rightarrow$	$f \ a \rightarrow$	$f \ b$

$(\langle \$ \rangle) ::$	$a \rightarrow$	$f \ b \rightarrow f \ a$
$(\langle * \rangle) ::$	$f \ a \rightarrow$	$f \ b \rightarrow f \ a$

Combinator zoo, expanded

$(\$)$	$::$	$(a \rightarrow b) \rightarrow$	$a \rightarrow$	b
$(<\$>)$	$::$	$(a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$
$(<*>)$	$::$	$f\ (a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$

$(<\$>)$	$::$	$a \rightarrow$	$f\ b \rightarrow f\ a$
$(<*>)$	$::$	$f\ a \rightarrow$	$f\ b \rightarrow f\ a$

Combinator zoo, expanded

$(\$)$	$::$	$(a \rightarrow b) \rightarrow$	$a \rightarrow$	b
$(<\$>)$	$::$	$(a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$
$(<*>)$	$::$	$f\ (a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$

const	$::$	$a \rightarrow$	$b \rightarrow$	a
$(<\$>)$	$::$	$a \rightarrow$	$f\ b \rightarrow$	$f\ a$
$(<*>)$	$::$	$f\ a \rightarrow$	$f\ b \rightarrow$	$f\ a$

Combinator zoo, classified

-- *Application-like*

$(\$)$	$::$	$(a \rightarrow b) \rightarrow$	$a \rightarrow$	b
$(<\$>)$	$::$	$(a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$
$(<*>)$	$::$	$f\ (a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$

-- *Const-like*

const	$::$	$a \rightarrow$	$b \rightarrow$	a
$(<\$>)$	$::$	$a \rightarrow$	$f\ b \rightarrow$	$f\ a$
$(<*>)$	$::$	$f\ a \rightarrow$	$f\ b \rightarrow$	$f\ a$

More combinators?

-- *Application-like*

$(\$)$::

$(<\$>)$::

$(<*>)$::

$(a \rightarrow b) \rightarrow$

$(a \rightarrow b) \rightarrow$

$f (a \rightarrow b) \rightarrow$

$a \rightarrow$

$f a \rightarrow$

$f a \rightarrow$

b

$f b$

$f b$

More combinators.

-- *Application-like*

$(\$)$::

$(<\$>)$::

$(<*>)$::

flap ::

$(a \rightarrow b) \rightarrow$

$(a \rightarrow b) \rightarrow$

$f (a \rightarrow b) \rightarrow$

$f (a \rightarrow b) \rightarrow$

$a \rightarrow$

$f a \rightarrow$

$f a \rightarrow$

$a \rightarrow$

b

$f b$

$f b$

$f b$

More combinators.

-- *Application-like*

$(\$)$	$::$	$(a \rightarrow b) \rightarrow$	$a \rightarrow$	b
$(<\$>)$	$::$	$(a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$
$(<*>)$	$::$	$f\ (a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$
$flap$	$::$	$f\ (a \rightarrow b) \rightarrow$	$a \rightarrow$	$f\ b$
$(=<<)$	$::$	$(a \rightarrow f\ b) \rightarrow$	$f\ a \rightarrow$	$f\ b$

Flipped combinators.

-- *Application-like*

$(\$)$	$::$	$(a \rightarrow b) \rightarrow$	$a \rightarrow$	b
$(<\$>)$	$::$	$(a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$
$(<*>)$	$::$	$f\ (a \rightarrow b) \rightarrow$	$f\ a \rightarrow$	$f\ b$
$flap$	$::$	$f\ (a \rightarrow b) \rightarrow$	$a \rightarrow$	$f\ b$
$(=<<)$	$::$	$(a \rightarrow f\ b) \rightarrow$	$f\ a \rightarrow$	$f\ b$

-- *Application-like, arguments flipped*

$(&)$	$::$	$a \rightarrow$	$(a \rightarrow b) \rightarrow$	b
$(<\&>)$	$::$	$f\ a \rightarrow$	$(a \rightarrow b) \rightarrow$	$f\ b$
$(<*>)$	$::$	$f\ a \rightarrow$	$f\ (a \rightarrow b) \rightarrow$	$f\ b$
$(??)$	$::$	$a \rightarrow$	$f\ (a \rightarrow b) \rightarrow$	$f\ b$
$(>>=)$	$::$	$f\ a \rightarrow$	$(a \rightarrow f\ b) \rightarrow$	$f\ b$

($\gg=$) allows more flexible parsing

Example: matrices

3	3	6	3	9	8	3	2	5	9	4	9	7	6
5	7	1	8	8	4	5	4	8		2	2	9	8
7	2	2	6	2	9	6	3	8	7	2			

✓ ✗ ✓ ✗ ✓

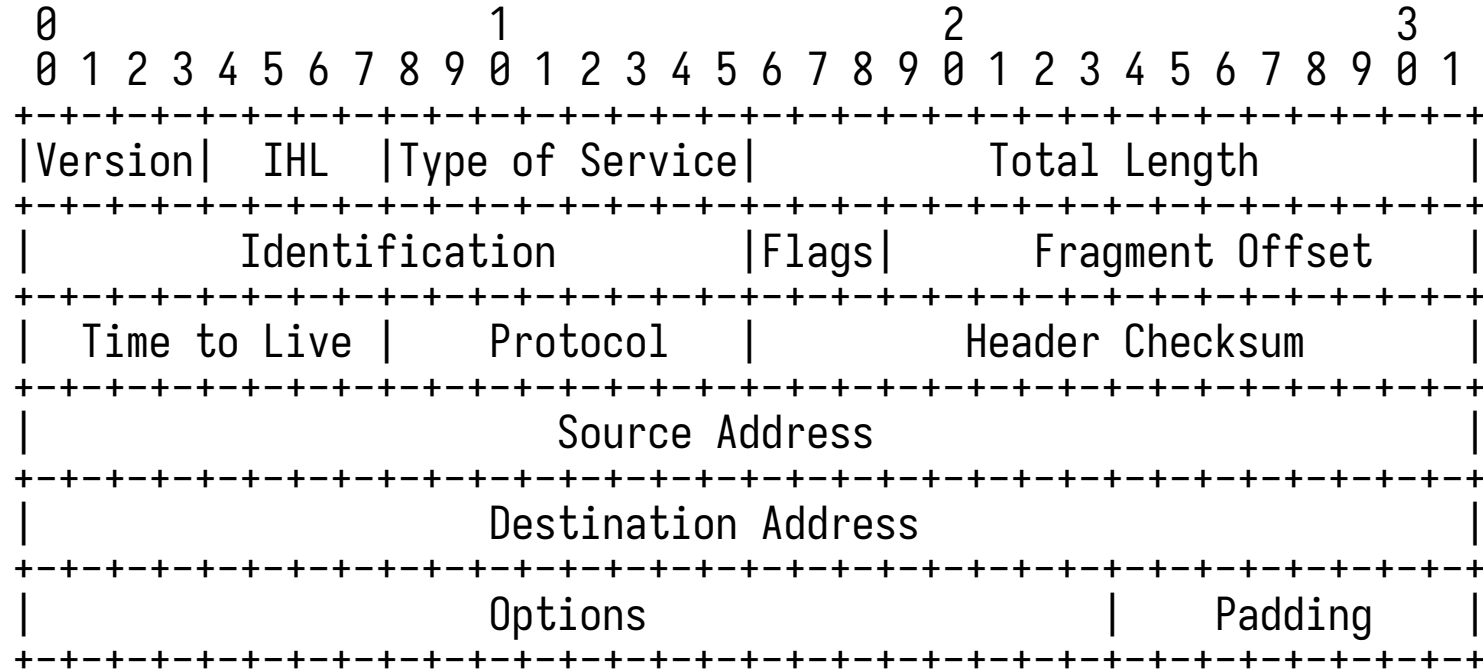
Numer of columns constant but not known in advance.

```
parseMatrix :: Parser Char [[Int]]
```

```
parseMatrix = parseRowOfAnyLength >>= \r →
```

```
  (r :) <$> parseRowsOfFixedLength (length r)
```

Example: IPv4 packets



Payload length not known in advance.

```
parseIpPacket : Parser Byte IpPacket
parseIpPacket = parseIpHeader >>= \h →
  (IpPacket h) <$> parseIpPayload (h .totalLength)
```

Example: nestable markdown divs

```
:::::::::::: columns
```

```

```

```
::::: bold
```

```
It quacks like a duck...
```

```
:::::
```

```
::::::::::::
```

```
parseMarkdown :: Int → Parser Line MarkdownBlock
```

```
parseMarkdown d = text <|> image <|> (parseDivHeaderShorterThan d >>= \h →  
  Div h <$> many (parseMarkdown (h .length))  
  <*> parseDivFooter (h .length) )
```

Haskell has special syntax for ($\gg=$)

-- You write

do

a ← foo

b ← bar

baz

-- Equivalent to

foo $\gg=$ (\a →

bar $\gg=$ (\b →

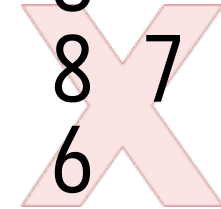
baz))

Quiz: grammar of well-formed matrices?

3 3
5 7
7 2



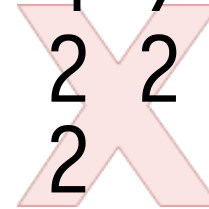
6 3
1 8
2 6



9 8 3 2
8 4 5 4
2 9 6 3



5 9 4 9
8 2 2
8 7 2



7 6
9 8



1. Matrix $\rightarrow$ Int**
2. Matrix $\rightarrow$ Int* Int**
3. Matrix $\rightarrow$ Int*
4. Matrix $\rightarrow$ Row Row*
Row $\rightarrow$ Int Int*

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Answer: none of the above!

- The matrix language is **not** context-free.
- ($\gg=$) can parse it anyway.
- Sometimes ($\gg=$) can be avoided, e.g.
 - `parseMatrix` $=$ many (many integer)
 - Check row-length later

Back to ($\gg=$)-free parsers

Quiz: parser for wait lines?

10:15 Utrecht 10:21

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$(,,) :: a \rightarrow b \rightarrow c \rightarrow (a,b,c)$

{- 1 -}	(,,)	<\$>	parseTime	<\$>	parseString	<\$>	parseTime
{- 2 -}	(,,)	<*>	parseTime	<\$>	parseString	<\$>	parseTime
{- 3 -}	(,,)	<*>	parseTime	<\$>	parseString	<\$>	parseTime
{- 4 -}	(,,)	<*>	parseTime	<\$	parseString	<\$	parseTime
{- 5 -}	(,,)	<\$>	parseTime	<*>	parseString	<*>	parseTime
{- 6 -}	(,,)	<\$>	parseTime	<*	parseString	<*	parseTime

“Applicative style”

| 10:15 Utrecht 10:21

$(,,) :: a \rightarrow b \rightarrow c \rightarrow (a,b,c)$

```
--      (,,)    <$> parseTime <$> parseString <$> parseTime
--      (,,)    <*> parseTime <$> parseString <$> parseTime
--      (,,)    <*> parseTime <$> parseString <$> parseTime
--      (,,)    <*> parseTime <$> parseString <$> parseTime
{- 5 -} (,,)    <$> parseTime <*> parseString <*> parseTime
--      (,,)    <$> parseTime <*> parseString <*> parseTime
```

Easy enough?

"10:15 Utrecht 10:21"

(,,) <\$> parseTime

<*> parseString

<*> parseTime

:: Parser Char (Time,String,Time)

Whitespace trouble

```
"10:15 Utrecht    10:21"
```

```
(,,) <$> parseTime
```

```
<*> parseString
```

```
<*> parseTime
```

```
:: Parser Char (Time,String,Time)
```

```
(AM 10 15, " Utrecht    ", AM 10 21)
```

Whitespace lexed away

```
["10:15", "Utrecht", "10:21"]
```

```
(,,) <$> parseTime
```

```
<*> parseString
```

```
<*> parseTime
```

```
:: Parser String (Time, String, Time)
```

```
(AM 10 15, "Utrecht", AM 10 21)
```

Lexing can be inconvenient

```
["10:15", "Den", "Haag", "10:21"]
```

```
(,,) <$> parseTime
```

```
<*> parseString
```

```
<*> parseTime
```

```
:: Parser String (Time, String, Time)
```

No parse.

Whitespace parsing

"10:15 Utrecht 10:21"

(,,) <\$> parseTime <*> some space
<*> parseString <*> some space
<*> parseTime

:: Parser Char (Time,String,Time)

(AM 10 15, "Utrecht", AM 10 21)

Demo time

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```
import ParseLib.Derived
```

```
time :: Parser Char Time
```

```
time = AM
```

```
    <$> natural <*> symbol ':'
```

```
    <*> natural
```

```
space = symbol ' '
```

```
string = many anySymbol
```

```
tline :: Parser Char
```

```
    (Time,String,Time)
```

```
tline = (,,)
```

```
    <$> time    <*> some space
```

```
    <*> string  <*> some space
```

```
    <*> time    <*> eof
```

Quiz: what does GHCi return?

```
ghci> parse tline "10:15 Utrecht    10:21"
```

Too many results!

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```
time :: Parser Char Time
time = AM
      <$> natural <*> symbol ':'
      <*> natural
```

```
space = symbol ' '
string = many anySymbol
```

```
tline :: Parser Char
         (Time,String,Time)
tline = (,,)
      <$> time    <*> some space
      <*> string  <*> some space
      <*> time    <*> eof
```

Quiz: what does GHCi return?

```
ghci> parse tline "10:15 Utrecht 10:21"
[((AM 10 15,"Utrecht ",AM 10 21),"")
,((AM 10 15,"Utrecht ",AM 10 21),"")
,((AM 10 15,"Utrecht",AM 10 21),"")]
```

Quiz: why are too many parses returned?

Definition of some

some :: Parser c a $\rightarrow$ Parser c [a]

some p = (:) <\$> p <*> many p

many :: Parser c a $\rightarrow$ Parser c [a]

many p = some p <|> succeed []

Expanding some

some space

== (:) <\$> space <*> many space

```

== (:) <$> space <*> (some space
    <|> succeed [])

```

```
= (:) <$> space <*> ((:) <$> space <*> many space
    <|> succeed [])
```

```
= (:) <$> space <*> ((:) <$> space <*> (some space
                                <|> succeed []))
    <|> succeed [])
```

Expanding some

some space

```
= (:) <$> space <*> ((:) <$> space <*> ((:) <$> space <*> some space))  
<|> (:) <$> space <*> ((:) <$> space <*> ((:) <$> space <*> succeed []))  
<|> (:) <$> space <*> ((:) <$> space <*> succeed [])  
<|> (:) <$> space <*> succeed []
```

<|> takes results from both parsers

some space

```
= (:) <$> space <*> ((:) <$> space <*> ((:) <$> space <*> some space))  
<|> (:) <$> space <*> ((:) <$> space <*> ((:) <$> space <*> succeed []))  
<|> (:) <$> space <*> ((:) <$> space <*> succeed [])  
<|> (:) <$> space <*> succeed []
```

$(\langle | \rangle) :: \text{Parser } s \ a \rightarrow \text{Parser } s \ a \rightarrow \text{Parser } s \ a$
 $(p1 \ \langle | \rangle \ p2) \ s = (p1 \ s) ++ (p2 \ s)$

But we only want the first result!

<<|> avoids calling second parser

some space

```
= (:) <$> space <*> ((:) <$> space <*> ((:) <$> space <*> some space))  
<|> (:) <$> space <*> ((:) <$> space <*> ((:) <$> space <*> succeed []))  
<|> (:) <$> space <*> ((:) <$> space <*> succeed [])  
<|> (:) <$> space <*> succeed []
```

$(\langle | \rangle) :: \text{Parser } s \ a \rightarrow \text{Parser } s \ a \rightarrow \text{Parser } s \ a$
 $(p1 \ \langle | \rangle \ p2) \ s = (p1 \ s) ++ (p2 \ s)$

$(\langle \langle | \rangle) :: \text{Parser } s \ a \rightarrow \text{Parser } s \ a \rightarrow \text{Parser } s \ a$
 $(p1 \ \langle \langle | \rangle \ p2) \ s = \text{if } (p1 \ s) \neq [] \text{ then } (p1 \ s) \text{ else } (p2 \ s)$

Definition of greedy, greedy1

greedy1 :: Parser c a → Parser c [a]

greedy1 p = (:) <\$> p <*> greedy p

greedy :: Parser c a → Parser c [a]

greedy p = greedy1 p <<|> succeed []

Quiz: which parsers always return exactly 1 result?

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- {- 1 -} many space
- {- 2 -} greedy space
- {- 3 -} (greedy space <* greedy space)
- {- 4 -} (greedy space <* many space)
- {- 5 -} (many space <* greedy space)
- {- 6 -} (many (greedy space))
- {- 7 -} (greedy (greedy space))
- {- 8 -} (greedy empty)
- {- 9 -} (many empty)

These parsers always return exactly 1 result

```
--      many    space
{- B -} greedy space
{- C -} (greedy space  < * greedy space)
{- D -} (greedy space  < * many space)
--      (many    space  < * greedy space)
--      (many (greedy space))
{- G -} (greedy (greedy space))
--      (greedy empty)
--      (many empty)
```

(many empty) recurses forever

(many empty)
== some empty <|> succeed []

(many empty) recurses forever

(many empty)

= some empty <|> succeed []

= ((:) <\$> empty <*> many empty) <|> succeed []

(many empty) recurses forever

(many empty)

= some empty <|> succeed []

= ((:) <\$> empty <*> many empty) <|> succeed []

= (map (() :) (many empty)) <|> succeed []

(many empty) recurses forever

(many empty)

= some empty <|> succeed []

= ((:) <\$> empty <*> many empty) <|> succeed []

= (map (():) (many empty)) <|> succeed []

= (map (():) (...)) <|> succeed []

(many empty) recurses forever

```
(many empty)
== some empty <|> succeed []
== ((:) <$> empty <*> many empty) <|> succeed []
== (map (() :) (many empty)) <|> succeed []
== (map (() :) (...)) <|> succeed []
```

- Laziness was not enough!

(many empty) recurses forever

```
(many empty)
== some empty <|> succeed []
== ((:) <$> empty <*> many empty) <|> succeed []
== (map (() :) (many empty)) <|> succeed []
== (map (() :) (...)) <|> succeed []
```

- Laziness was not enough!
- Beware left-recursion.

Summary (today)

- $(\gg=)$, $(\langle * \rangle)$, $(\langle \$ \rangle)$, $(\$)$
- $(\langle * \rangle)$, $(\langle \$ \rangle)$, **const**
- greedy, many
- Excess results
- Left-recursion

How to build parsers

1. Design grammar
2. Write parser
3. Discover bugs
4. Fix parser, transform grammar
5. GOTO 3